

# Rigidity percolation

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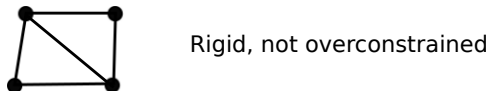
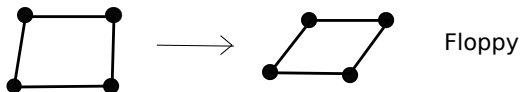
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# Rigidity

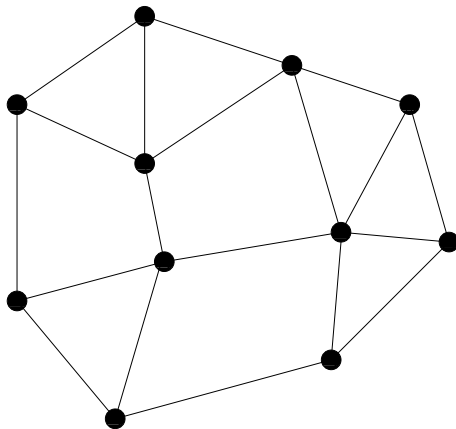
- Consider bars, which have a fixed length, linked together by "joints". Is the system **rigid** or **floppy** ?

Example in 2 dimensions; bar lengths are fixed, not the angles:



# Rigidity

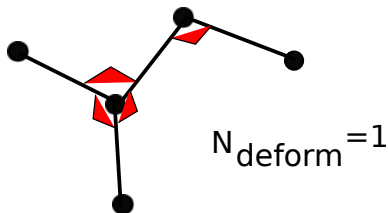
- When there are only a few joints and bars, it is easy...  
What about this network, with 11 sites?



- Is it floppy? Rigid? How many floppy modes? Where?

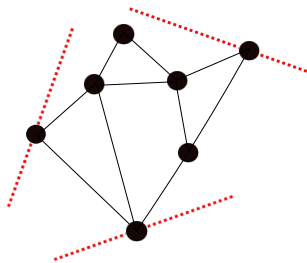
## Related problems

- Bond bending constraints: angles between two adjacent bonds have to be kept fixed ( $D = 3$ )



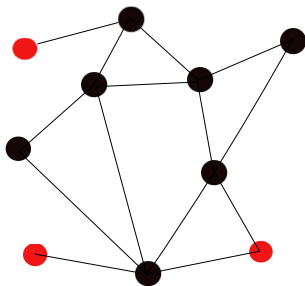
## Related problems

- Bond bending constraints: angles between two adjacent bonds have to be kept fixed ( $D = 3$ )
- Rigidity with "sliders": some joints constrained to move on a line



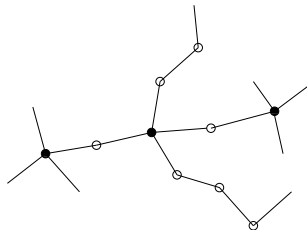
## Related problems

- Bond bending constraints: angles between two adjacent bonds have to be kept fixed ( $D = 3$ )
- Rigidity with "sliders": some joints constrained to move on a line
- Rigidity with "pinned" joints, which cannot move at all



## An application : "covalent glasses"

- Example: a disordered network with Germanium and Selenium atoms. Ge = 4 bonds; Se = 2 bonds.
- Bond lengths and angles between two adjacent bonds can be considered as constraints ( $\sim$  the energy needed to modify them is larger than the temperature).

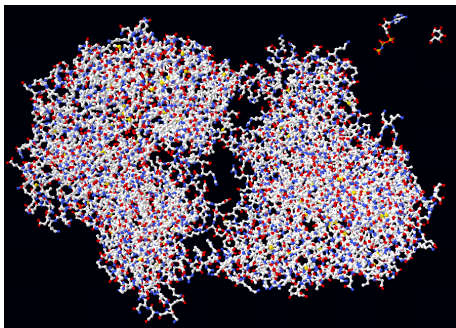


Each bond: 1 length constraint; each Se atom: 1 angular constraint; each Ge atom: 5 angular constraints.

→ Go from "floppy" to "rigid" by increasing the Ge fraction.

## Another application: protein rigidity (MF Thorpe and coworkers)

- Proteins are large biological molecules. An example (hexokinase):



Let's simplify:

Atoms → balls; chemical (or other strong) bonds → bonds; weak interactions → forgotten!

→ is the simplified structure floppy or rigid?

→ if floppy, what are the possible deformations?



# Constraint counting

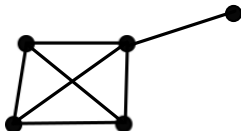
**Maxwell's idea:** constraint counting

- each joint starts with 2 degrees of freedom
  - each bar removes one degree of freedom
- First try: formula for the number of remaining degrees of freedom,  $N_{d.o.f.}$ ;  $N$  joints,  $M$  bars:

$$N_{d.o.f.} = 2N - M \text{ if } M < 2N - 3 ; N_{d.o.f.} = 3 \text{ if } M \geq 2N - 3$$

- Cannot be correct... Need to count redundant constraints:

$$N_{d.o.f.} = 2N - M + N_{\text{redundant}}$$



$$N=5; M=7$$

$$N_{\text{redundant}} = 1$$

$$N_{d.o.f.} = 4$$

## From geometry to graph theory: Laman theorem

- Power of constraint counting: replace a geometrical problem by a discrete, graph theoretical one.

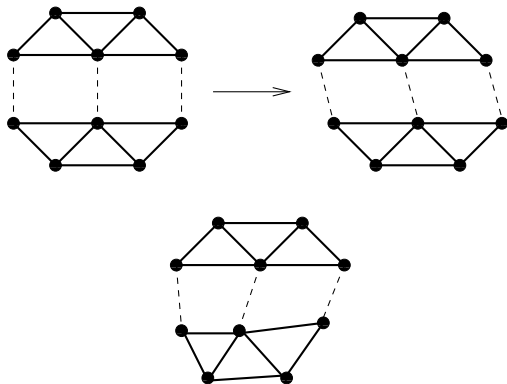
**Question:** is it possible to keep this desirable feature, correcting the approximations of constraint counting?

- **Generic rigidity** in 2D can be characterized in a purely graph theoretical way (Laman 1970):

$G$  has a redundant constraint  $\iff$  there is a subgraph with  $n$  vertices,  $m$  edges and  $m > 2n - 3$ .

$\rightarrow$   $\sim$  constraint counting on each subgraph to detect redundant constraints

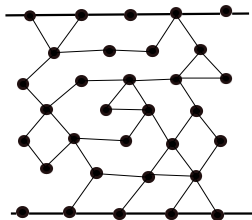
## Generic rigidity



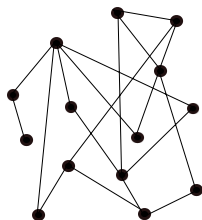
Top: a non generic realization; Laman theorem does not apply.  
Bottom: a generic realization of the same graph.

## Second ingredient: probabilities

In many cases, the structure is too large to be known exactly (think of covalent glasses for instance) → one would like to use a probabilistic description



Each link between a pair of neighboring vertices is present with proba.  $p < 1$

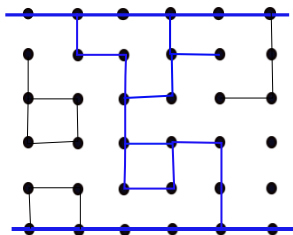


Links put "randomly", no geometry.

It is a **percolation** problem.

# "Standard" percolation

- "connectivity" percolation = well studied since the 60's



- Each link is present with proba.  $p$ , and absent with proba.  $1-p$

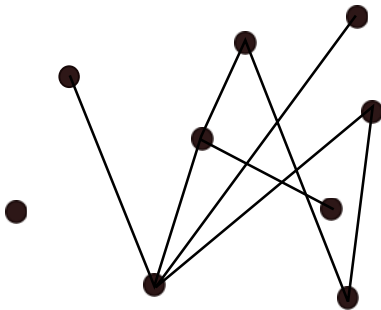
- Question: is there a path from top to bottom?

NB1: standard percolation is analog to "rigidity" percolation with one "degree of freedom" per vertex

NB2: standard percolation on a random graph = appearance of a "giant connected component"

# Erdos-Renyi random graphs

Definition of  $\mathcal{G}(n, c/n)$ :  $n$  vertices; any pair of vertices connected with proba.  $c/n$ . There is no notion of space.



Some properties: approximately  $nc/2$  edges; Poisson  $\mathcal{P}(c)$  degree distribution; **few small loops**...

## Questions for rigidity percolation

- Is there a well defined threshold  $p_c$  for the appearance of a "macroscopic rigid cluster"?

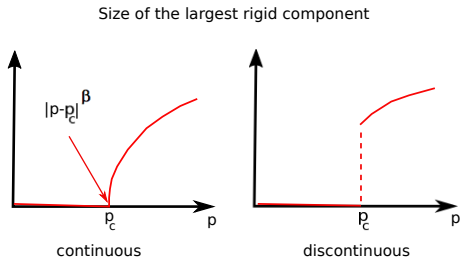
$$p < p_c \Rightarrow \text{percolation probability} = 0$$

$$p > p_c \Rightarrow \text{percolation probability} = 1$$

**Answer:** **yes** for random graphs and lattices (Numerics in the 90's; Holroyd  $\sim$ 2000); threshold computed by Kasiwisvanathan, Moore and Theran (KMT 2011) for  $\mathcal{G}(n, c/n)$  random graphs, unknown for lattices.

# Questions for rigidity percolation

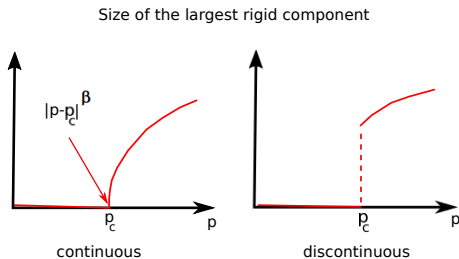
- Size of the largest rigid component? Continuous/discontinuous at  $p_c$ ?





## Questions for rigidity percolation

- Size of the largest rigid component? Continuous/discontinuous at  $p_c$ ?



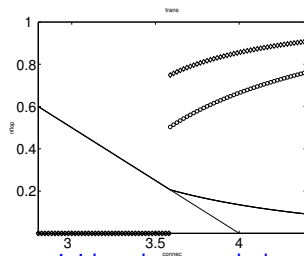
### Answer:

- Discontinuous for  $\mathcal{G}(n, c/n)$  random graphs (Theran)
- seems to be continuous for lattices (Jacobs-Thorpe, Duxbury-Moukarzel 90's, numerics).

## Questions for rigidity percolation

- Size of the largest rigid component? Continuous/discontinuous at  $p_c$ ?

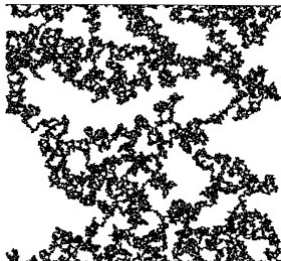
Example: Erdős-Rényi random graph  $\mathcal{G}(n, c/n)$ . Vary  $c$



Size of the biggest rigid and stressed clusters, and number of "floppy modes" vs mean connectivity

## Questions for rigidity percolation

- For lattices, what happens close to threshold? = "Critical" behavior?  $\beta = ?$  (critical exponent, exciting for statistical physicists); **fractal dimension?**



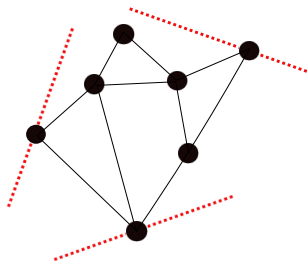
Overconstrained regions (Simulation by P. Duxbury et al.)

**Answer:** **unknown.** Critical exponents seem to be different from standard percolation.

# Goals

Fully understand the 2D lattice case: difficult... More modest goals:

1. Find models that can be solved;
  2. Explore similarities/differences standard percolation/rigidity percolation: study models that interpolate between both.
- Study rigidity percolation with sliders on random graphs



→ Study other kind of "simple" lattices (eg. hierarchical).

## Rigidity with sliders

- Consider a structure with  $n_1$  sites with sliders,  $n_2$  free sites and  $m$  bars. One slider = one constraint  
→ modify constraint counting

Difficulty: sliders "pin" the rigid components to the plane

- Distinguish between free, partly pinned, and pinned rigid clusters

**A Laman-type theorem** (I. Streinu, L. Theran, 2010).

Redundant constraint  $\iff$  subgraph with

$$n'_1 + 2n'_2 - m' - \max(3 - n'_1, 0) < 0$$

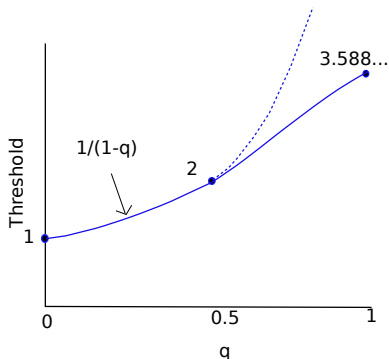
- A graph theoretical approach possible (under a genericity condition, as usual)

# Rigidity percolation with sliders

- Erdős-Renyi random graph  $\mathcal{G}(n, c/n)$ , with  $n = n_1 + n_2$   
 $n_1 = (1 - q)n$ ,  $n_2 = qn$ .  
 $1 - q =$  proportion of sites with sliders
- $q = 0$ : ordinary percolation = well known; continuous
- $q = 1$ : rigidity percolation, discontinuous; threshold  $c = 3.588\dots$
- **What happens in between?**

# Threshold

- percolation threshold vs proportion of sliders



- $c^* = 1/(1 - q)$  for  $q \leq 1/2$
- For  $q > 1/2$ , implicit expression for  $c^*(q)$ :

$$c^* = \frac{\xi^*}{1 - e^{-\xi^*} - q\xi^*e^{-\xi^*}}, \frac{\xi^*(1 - e^{-\xi^*} - q\xi^*e^{-\xi^*})}{(1 + q)(1 - e^{-\xi^*} - q\xi^*e^{-\xi^*}) - q(\xi^*)^2e^{-\xi^*}} = 2.$$

## Rigidity percolation with sliders, 2

**Theorem:** (JB, M. Lelarge, D. Mitsche)

Let  $G \sim \mathcal{G}(n, c/n)$  an Erdos-Renyi random graph, with a fraction  $1 - q$  of sliders. Then, we can compute  $c^*(q)$ , such that with high probability (proba  $\rightarrow 1$  when  $n \rightarrow \infty$ ):

- ▶ If  $c < c^*(q)$ , there is no giant rigid component
- ▶ If  $c > c^*(q)$ , there is a giant rigid component

Furthermore, for  $q < 1/2$  the transition is continuous, and for  $q > 1/2$  it is discontinuous.

**NB:**  $c^*(q = 0) = 1$  and  $c^*(q = 1) = 3.588 \dots$



## Size of the largest rigid component

- Size of the largest component at threshold: jump for  $q > 1/2$  :  
~ rigidity without sliders.
  - Continuous transition for  $q < 1/2$ : ~ connectivity percolation.
  - Discontinuous transition for  $q > 1/2$ .
- *tricritical point* at  $q = 1/2$  (statistical mechanics jargon)

## Strategy of proof

- **Step 1:** Link with *orientability* (generalizes the case without sliders)

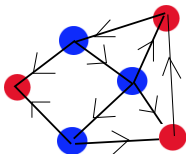
-Intuition: one bond removes one degree of freedom to one of the two vertices it links

-Vertices with or without slider: 1 or 2 degree of freedom

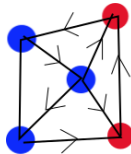
→ Link with "orientability"

● : with slider

● : free vertex



orientable



not orientable

## Strategy of proof, 2

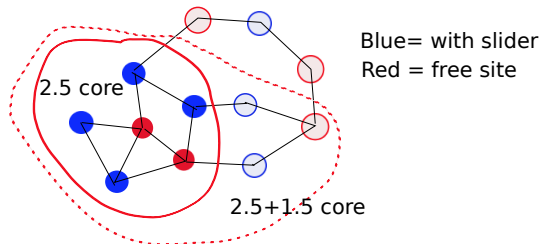
- **Step 2:** Thresholds for orientability and percolation are equal  
"Rigid"  $\Rightarrow$  "Non orientable" = easy  
"Non orientable"  $\Rightarrow$  "Rigid" = more laborious
- **Step 3:** Compute the threshold for orientability  $\rightarrow$  method introduced by M. Lelarge  
 $\sim$  rigorous "cavity method", a heuristic introduced by physicists.

## Step 4: type of transition

- For  $q > 1/2$  ("rigidity-like" transition), a density argument applies: rigid components must be dense enough, and dense subgraphs must have a minimal size of order  $n$  (uses again the generalization of L. Theran's lemma).

→ discontinuous transition

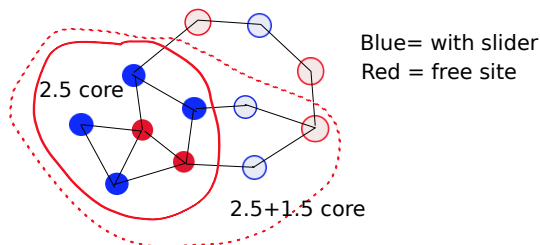
- For  $q < 1/2$  ("connectivity-like" transition), we need "cores"



Remove recursively blue sites with less than 2 links and red sites with less than 3. What remains is the "2.5-core". Then add recursively blue sites with one link to the core, and red sites with 2. One gets the "2.5 + 1.5-core".

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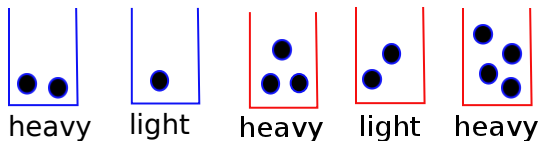
Then show: largest rigid component  $\subset 2.5 + 1.5$ -core

- Compute the size of the  $2.5 + 1.5$ -core and show it is small.

## Step 5: Size of cores

- Size of the  $3 + 2$  core = conjecture in Kasivisvanathan-Moore-Theran 2011.
- Strategy: use Janson-Luczak technique

Bins = vertices, with sliders (blue) or without (red)  
Balls = half edges



- good knowledge of degree distributions after the core construction
- possible to control the process growing the  $3+2$  core.

## Conclusions on random graphs

- ▶ Complete phase diagram with a tricritical point
- ▶ Proof combines many "old" ideas: strategy Theran et al. relating to orientability; M. Lelarge's technique to compute orientability threshold; Janson-Luczak technique to compute the size of cores
- ▶ What about rigidity with some pinned sites? Conjecture by physicists (Moukarzel '03): the discontinuous transition may disappear, but there is no continuous transition. . . *A proof seems accessible -joint work with Dieter Mitsche and Louis Theran*
- ▶ Physics literature: tree-like heuristics give access to much more detailed results (Large Deviation Cavity Method); could these be transformed into theorems? A general question, beyond rigidity.

## Beyond random graphs?

- ▶ Random graphs: much easier than percolation problems on lattices . . .
- ▶ whereas problems on lattices, or at least on graphs with some geometric content, are a priori more interesting for physics.
- ▶ Understand the phase transition on regular lattices (beyond existence proof by Holroyd)? Precise numerical simulations would be useful; I don't even have heuristic theoretical ideas...  
→ a lot to do here!